

# DNA

## Data and Program Representation

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Alexandre David  
1.2.05  
adavid@cs.aau.dk





# Introduction

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- Very important to understand how data is represented.
  - operations
  - limits
  - precision
- Digital logic built on 2-valued logic system
  - high/low 5V/0V true/false
  - we abstract from that from now on  
→ bits 0/1



# Basics

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## Natural/Real Numbers

- Base 10
- Infinite
- Exact

## Computer Numbers

- Base 2
- Finite
- Rounding - overflow

## In this lecture

- How to represent numbers & characters – range, encoding.
- A little arithmetic.
- How to use these numbers.



# Questions

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- How to code negative numbers?
- How to code real numbers?
- Which kind of precision do we get?
  - Small numbers vs. big numbers.
- What about characters?



# Example

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- Overflow:

```
main() {  
    printf("%d\n",200*300*400*500);  
}
```

outputs -88490188.

- Fix – sort of:

```
main() {  
    printf("%lld\n",200LL*300LL*400LL*500LL);  
}
```

outputs 120000000000.

What if I forget LL?



# Example

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- Loss of precision:  
 $(3.14+1e20)-1e20==0.0$   
 $3.14+(1e20-1e20)==3.14$
- Test  $x == 0.0$  not very useful when solving equations.
- In this lecture you will know why.



# Data Storage

- Basic unit is the byte (= 8 bits).

C-declaration	Typical 32-bit	Typical 64-bit
char	1	1
short int	2	2
int	4	4
long int	4	8
char *	4	8
float	4	4
double	8	8



# “Features”

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- Limits on addressable memory.
- Size linked to architecture – 32/64.
- Aligned memory allocation (32/64 bits).
- Careful on addressing:

```
main() {  
    char a[]="Hello world!";  
    int *p=&a[1];  
    printf("%d\n",*p);  
}
```

Bus error on some CPUs





# Integer Encoding

- Unsigned integers:

$$UB = \sum_{i=0}^{w-1} x_i 2^i$$

- Signed integers:  
Called **2 complement**.

$$SB = -x_{w-1} 2^{w-1} + \sum_{i=0}^{w-2} x_i 2^i$$

w: size of a word (in bits) x: bits (0 or 1)
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- Highest bit codes the sign.



# Range & Examples

- Examples:  $1010 = 10$ ,  $0110 = 6$ ,  $0101 = 5$

$2^3=8$	$2^2=4$	$2^1=2$	$2^0=1$
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- Range:
  - unsigned  $2^k$  numbers from 0 to  $2^k-1$
  - signed  $2^k$  numbers from  $-2^{k-1}$  to  $2^{k-1}-1$ 
    - one more negative number than positive ones
- How to convert between types?
  - int – char – long int...
  - sign extension



# Basic Arithmetic

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- Logical operations (bitwise):  
 $\&, |, ^, \sim, \ll, \gg$ .  
Example:  $a \wedge b; b \wedge a; a \wedge b;$
- Arithmetic operations:  $+ - * /$ .
- Careful with shifts on signed integers!
  - arithmetic & logical shifts
- Do not mess up with boolean operators ( $\&\&, ||$ ).



# Properties

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- Most operators are the same on signed/unsigned integers – from a binary point of view – beauty of the encoding.
  - One hardware implementation for  $+ - / * \dots$  valid for signed and unsigned integers.
  - Example on 4 bits:  
 $1 + 1001 = 1010$   
unsigned:  $1+9 = 10$   
signed:  $1+(-7) = -6$



# Properties

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- Operations based on the algebra  $\langle \mathbb{Z}_n, +_n, *_n, -_n, 0, 1 \rangle$  (commutativity, associativity, distributivity, ...)  
Operations modulo  $n$  and  $-a=0$  or  $-a=n-a$ .
- Similar to boolean algebra  $\langle \{0,1\}, |, \&, \sim, 0, 1 \rangle$  with the addition of DeMorgan laws  
 $\sim(a \& b) = \sim a | \sim b$ ,  $\sim(a | b) = \sim a \& \sim b$



# Practice: Shifts & Masks

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- Read bit  $n$ : Use mask  $(1 \ll n)$ .
- Set bit  $n$  on int `bits[]`:  
`ipos = n / 32;`  
`imask = 1 << (n % 32);`  
`bits[ipos] |= imask;`
- Division/multiplications by powers of 2 seen as shifts.
- 2-complement:  $-a = \sim(a-1) = \sim a + 1$



# Arithmetic


- Machine code of  $+$   $-$   $*$   $/$  same for int/uint.
- Integer conversion == type casting.
  - Padding for the sign (int).
  - Conversion is modulo the size of the new int.
  - Beware of implicit conversions in C!
- Optimizations for some operations:

```
2*a      == a+a  == a<<1      a/2      == a>>1
a*2^i    == x << i      a/2^i    == a>>i
a%2^i    == a & ((1<<i)-1)  2^i    == 1<<i
```



# Notes

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- Beware of precedence of operators:
  - if (x & mask == value) WRONG
  - if ((x & mask) == value) RIGHT
- Test odd numbers: if (x & 1)
- Careful:  
unsigned int i;  
for(i = 0; i < n-1; ++i) ... 





# Overflow – “carry”

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$$\begin{array}{r}
 1011 \\
 +1101 \\
 \hline
 111 \\
 \hline
 11000
 \end{array}$$

Subtraction?

1011	multiplicand
*1101	multiplier
<hr/>	
1011	
0000	partial products
1011	
1011	
<hr/>	
10001111	product



# Hexadecimal Notation

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- Learn the first powers of 2.
- Hexadecimal more useful:
  - One digit codes 4 bits.  $0...F=0...15=16$  numbers.
  - Notation:  $0x...$
  - Why?
- Examples:
  - $0xa57e = 1010\ 0101\ 0111\ 1110$   
 $10=8+2, 5=4+1, 7=4+2+1, 14=8+4+2$
  - $0xf = 1111, 0x7=0111, 0x3=0011$



# Endianness: Beware!

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- “0xa57e” is a notation for humans.  
Corresponds to “1010 0101 0111 1110” in base 2.
  - Little endian: stored as 0111111010100101.
  - Big endian: stored as 1010010101111110.
- Does not matter in C/Java/C#, except for
  - bitmap manipulation
  - device drivers
  - network transfers



# Testing for Endianness

- Write a value on 32 bits.
  - Read 8/16 bits and check what was written.
  - Exercise for Sun/Intel.

```
main() {  
    int a = 0xf0000000;  
    char *c = &a;  
    printf("%x\n", *c);  
}
```

- Intel? → 0
- Sun? → 0xffffffff0
- What if a = 0x70000000?



# Representation of Reals

- How to code a real number with bits?
  - Finite precision → approximation.
  - Represent very small and very large numbers → **density** of encoding varies.
- Scientific notation used, e.g. (base 10), 3.141e2 – but in base 2.
- Starter: fractional numbers – bad for large or small numbers.

- Decimal (d):

- Binary(b):

$$d = \sum_{i=-n}^m 10^i d_i \quad b = \sum_{i=-n}^m 2^i b_i$$



# IEEE Floats

- IEEE 754 floating point standard



- $V = (-1)^S M \cdot 2^E$
- Number of bits (float/double):  
s[1], m[23/52], e[8/11].
- Normalized and de-normalized values.
- Bit fields: s, m, e to code respectively S, M, E.

- Normalized values ( $e \neq 0, e \neq 111\dots$ )

- $E = e - \text{bias}$  (-126...127/-1022...1023).
- $M = 1 + m$  ( $1 \leq M < 2$ )
- Trick for more precision: implied leading 1.



# IEEE Floats

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- De-normalized values ( $e=0$  or  $11\dots$ )
  - $e=0$ :  $E=1$ -bias,  $\text{bias}=2^{k-1}-1$ 
    - Coding compensates for  $M$  not having an implied leading 1.
    - $M=m$
    - For numbers very close to 0.
  - $e=11\dots$ :
    - $m=0$ , (signed infinite)
    - $m\neq 0$ , NaN.

No more of this.  
We abstract from  
this complexity.



# Ranges of Floats

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- Single precision (float)  
 $2^{-126} \dots 2^{127} \sim 10^{-38} \dots 10^{38}$ .
- Double precision (double)  
 $2^{-1022} \dots 2^{1023} \sim 10^{-308} \dots 10^{308}$ .





# “Features” of IEEE floats

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- $+0.0 == 0$  (binary representation = 00...).
- If interpreted as unsigned int, floats can be sorted (+x ascending, -x descending).
- All int values representable by doubles.
- Not all int values representable by floats.
  - round to even (avoid stat. bias)
  - round towards 0
  - round up
  - round down
  - can't choose in C...



# Properties (floats)

- Operations **NOT** associative.
- Not always inverse (infinity).
- Loss of precision.
- Ex:  $x = a + b + c$ ;  $y = b + c + d$ ;  
Optimize or not?

Important for  
compilers and  
programmers.

- Monotonicity  $a \geq b \Rightarrow a + x \geq b + x$
- Casts:
  - int2float rounded, double2float rounded/overflow
  - int2double, float2double OK
  - float2int, double2int truncated/rounded/overflow.



# IA32: The Good And The Bad

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- Good: Uses internally 80 bits extended registers for more precision.
- Bad:
  - Stack based.
  - Side effects like changing values when loading or saving numbers in memory whereas register transfers are exact.
- Extensions: MMX, SSE, (AltiVec). SIMD instructions = operations working in parallel on multiple data.



# Numerical Precision

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- Evaluation of precision
  - absolute  $x \pm \alpha$
  - relative  $x*(1 \pm \alpha)$
- Be careful with division by very small values:  
Can amplify numerical errors.
  - Numerical justification for Gauss' method to solve linear equations.



# Character Sets

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- By convention (standard), we assign codes to characters.
  - Careful with programming languages  
C char = 1 byte, Java char = 2 bytes
  - ASCII (American Standard Code for Information Interchange) common (7 bits), extended ASCII (8 bits).
  - Unicode
  - Other IEEE8859-x codings.
  - Other Japanese codings...



# More Complex Structures

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- How to represent, e.g.,  
struct { int a, b; char c}?
- Use continuous bits for the successive fields.
  - Compilers will *align* data!
  - Try:

```
typedef struct { int a; char c; } foo_t;
int main()
{
    printf("%d\n", sizeof(foo_t));
    return 0;
}
```



# What About Programs?

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- Instructions coded into bytes.
  - “opcode”
  - Processors interpret them as their particular instruction sets (standard).