Process-Processor Mapping (2.7)

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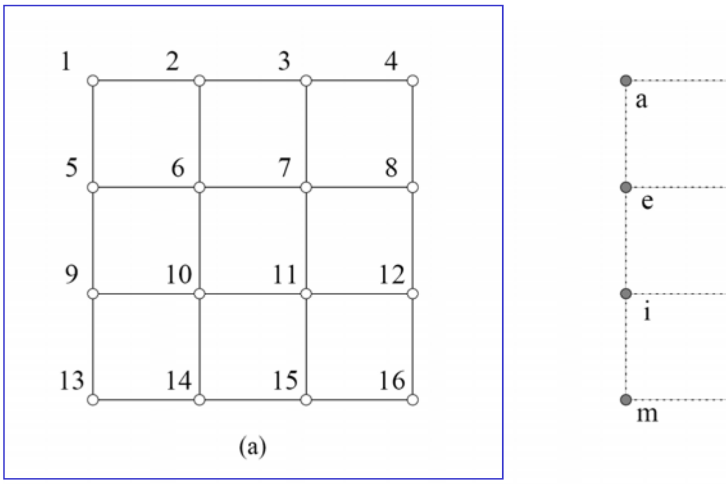
1.2.05

Example

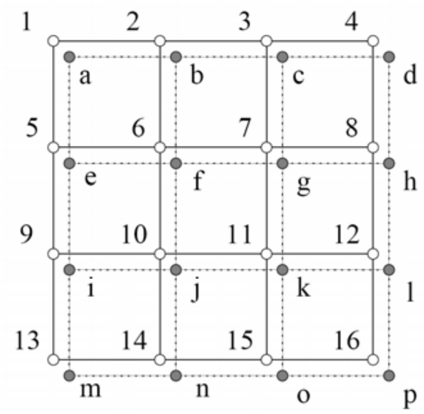
Underlying architecture (physical network).
Processors.

Processes and their interactions.



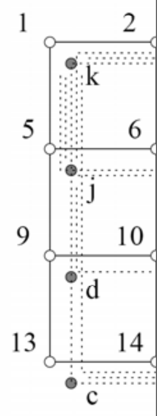
Example

Intuitive mapping.



(c)

Random mapping and congestion.



Here we have congestion because of the mapping although the intuitive mapping didn't have it.

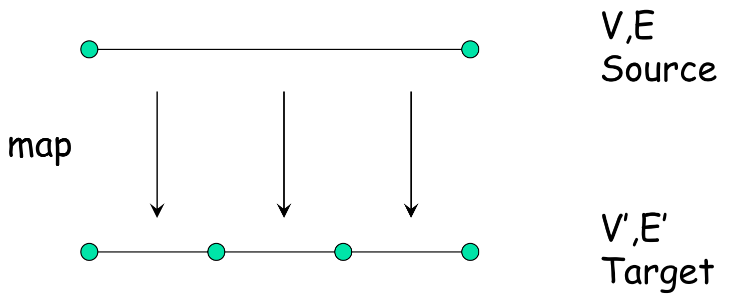
Mapping Techniques For Graphs

- Topology embedding:
  - Embed a communication pattern into a given interconnection topology.
    Hypercube in a 2-D mesh?
    2-D mesh in a hypercube?
- Why?
  - Cost.
  - Design an algorithm for a topology but you port it to another.

Embedding Metrics

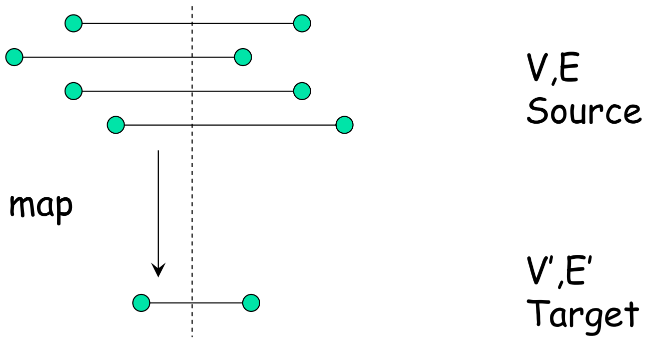
- Map a graph *G(V,E)* into *G′(V′,E′)*.
  - Dilation: Maximum number of links of E′ an edge of E is mapped onto.
  - Expansion: ratio |V′|/|V|.
  - Congestion: Maximum number of edges of E mapped on a single link of E′.

Dilation & Expansion



Dilation: 3.
Expansion: 4/2 = 2.

Congestion



Congestion: 4.

Typo fig 2.29, congestion is 5, not 6.

Embedding a Linear Array Into a Hypercube

- Map a linear array (or ring) of 2^d nodes into a d-dimensional hypercube.
- How would you do it? ?
- Gray code function:

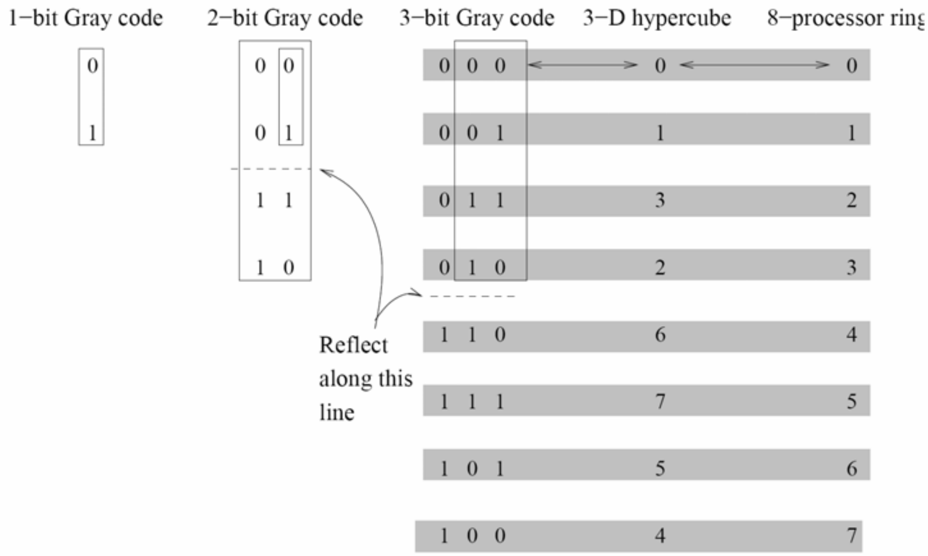
$$G(0,1) = 0$$

$$G(1,1) = 1$$

$$G(i, x+1) = \begin{cases} G(i,x), & i < 2^x \\ 2^x + G(2^{x+1} - 1 - i, x), & i \geq 2^x \end{cases}$$

Where to map node i into hypercube of dim d: G(i,d).

Gray Code



1−bit Gray code
2−bit Gray code
3−bit Gray code
3−D hypercube
8−processor ring
0
1
0 0
0 1
1 1
1 0
0 0 0 0 0
0 0 1 1 1
0 1 1 3 2
0 1 0 2 3
1 1 0 6 4
1 1 1 7 5
1 0 1 5 6
1 0 0 4 7
Reflect
along this
line

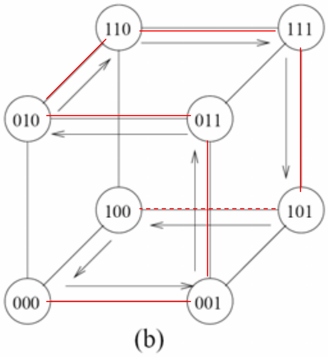


Figure 2.30 (a) A three-bit reflected Gray code ring; and (b) its embedding into a three-dimensional hypercube.

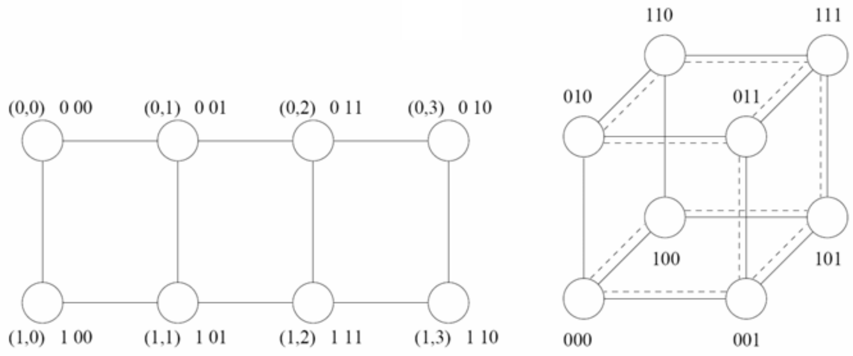
Gray Code Mapping

Gray Code Mapping cont.

- *G(i,d)* : i^{th} entry in sequence of *d* bits.
- Adjoining entries *G(i,d)* and *G(i+1,d)* differ from **only one bit**.
  - Like hypercubes -> direct link for these nodes.
- ? Dilation?
- ? Congestion?

Embedding a Mesh into a Hypercube

- Map a $2^r \times 2^s$ wraparound mesh into a r+s dimension hypercube.
- ? How?
- Map *(i,j)* to *G(i,r)||G(j,s)*. *Typo: no -1, it doesn't work.*
  - Extension of previous coding.



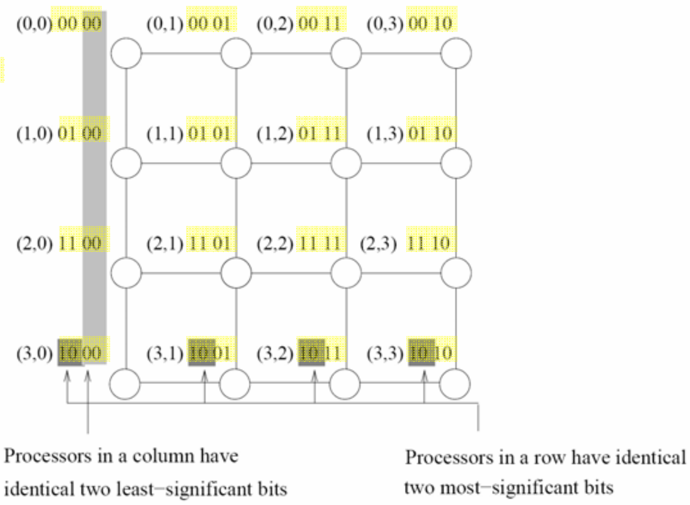
2x4 mesh into a 3-D hypercube

Embedding a Mesh into a Hypercube

- Properties
  - Dilation & congestion 1 as before.
  - All nodes in the **same row** (mesh) are mapped to hypercube nodes with ***r*** **identical most significant bits**.
  - Similarly for columns: s identical least significant bits.
  - What it means: They are mapped on a sub-cube!

Sub-Cube Property (4x4)

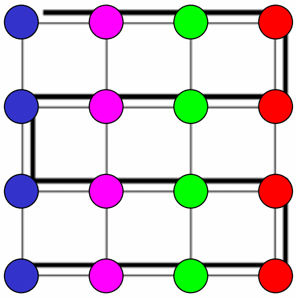
Gray codes



Embedding of a Mesh Into a Linear Array

- This time denser into sparser.
- 2-D mesh has $2p$ links and an array has p links.
  - There must be congestion!
  - Optimal mapping: in terms of congestion.

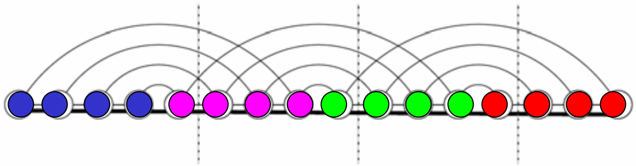
Easy: Linear Array Into Mesh



(a) Mapping a linear array into a
2D mesh (congestion 1).

Mesh Into Linear Array

Congestion: 5.



(b) Inverting the mapping − mapping a 2D mesh into a linear array (congestion 5)

Figure 2.32 (a) Embedding a 16 node linear array into a 2-D mesh; and (b) the inverse of the mapping. Solid lines correspond to links in the linear array and normal lines to links in the mesh.

What does it mean?

Is It Optimal?

- Bisection of
  - 2-D mesh is sqrt(p).
  - linear array is 1.
- 2-D mesh mapped on linear array has congestion *r*.
  - Cut in half linear array: cut 1 link, but cut no more than *r* mapped mesh links.
  - Lower bound: $r \geq sqrt(p)$.

The congestion has the lower bound given by bisection width of the original topology divided by the bisection width of the target topology.

•2D mesh → linear array: sqrt(p).

•2D mesh → ring: sqrt(p)/2.

•Hypercube → 2D mesh: (p/2)/sqrt(p) = sqrt(p)/2.

•Hypercube → wrap around 2D mesh: sqrt(p)/4.

Hypercube Into a 2-D Mesh

- Denser into sparser again (in terms of links).
- *p* even power of 2.
- *d*= log *p* dimension.
- d/2 least (most) significant bits define sub-cubes of $\sqrt{p}$ nodes.
- Row/column ↔ sub-cube, inverse of hypercube to 2-D mesh mapping.

$p=2^d$, d even.

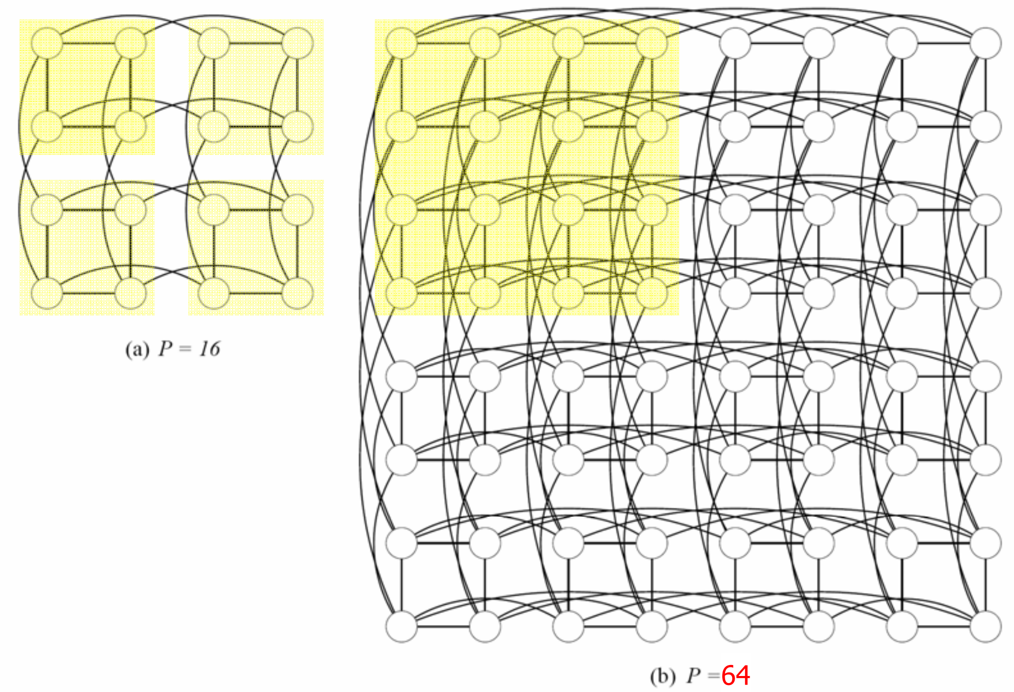


Figure 2.33 Embedding a hypercube into a 2-D mesh.

What Is The Point?

- Possible to map denser into sparser:
  - Map (expensive) logical topology into (cheaper) physical hardware!
  - Mesh with links faster by $\sqrt{p}/2$ than hypercube links has same performance!

Cost-Performance

- Read 2.7.2.
- Remember that 2-D mesh is better in terms of performance/cost.

Don’t be confused:

Wrap mesh sqrt(p)*sqrt(p) nodes, 4p/2 channels.

P nodes hypercube dim log(p), p*dim/2 wires = p*log(p)/2.