

Pthread Synchronization Patterns

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Creating + Running a Thread

Main program	Thread function
<p>pthread_create</p> <p>...</p> <p>pthread_join</p>	<p>...</p> <p>pthread_exit</p>

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Mutex-Lock

Main program	Thread
<p>pthread_mutex_init</p> <p>...</p> <p>pthread_mutex_destroy</p>	<p>pthread_mutex_lock</p> <p style="color: red; font-weight: bold;"><critical section as small as possible></p> <p>pthread_mutex_unlock</p>

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Mutex-Try-Lock

Main program	Thread
<pre>pthread_mutex_init ... pthread_mutex_destroy</pre>	<pre>if pthread_mutex_trylock <postpone work do something else> else <critical section as small as possible> pthread_mutex_unlock end</pre>

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Condition Variables & Monitors

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Condition Variables

Main program	Thread: Wait for a condition.
<pre>pthread_mutex_init pthread_cond_init ... pthread_mutex_destroy pthread_cond_destroy</pre>	<pre>pthread_mutex_lock while !condition do pthread_cond_wait done <critical section as small as possible> pthread_mutex_unlock</pre>
<p>Thread: Make a condition become true.</p>	<pre>pthread_mutex_lock <critical section as small as possible> pthread_cond_signal pthread_mutex_unlock</pre>

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Spin-Locks

- Mutex:
 - Threads block until the lock is acquired.
 - Blocked threads are idle and need to wake up.
- Spin-locks:
 - Threads spin until the lock is acquired.
 - Blocked threads are not idle!
 - Better for quick access of **small** critical sections with low contention.

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Pthread Spin Locks

- Calls:
 - `pthread_spin_init(pthread_spinlock_t*, int)`
`pthread_spin_destroy(pthread_spinlock_t*)`
 - `pthread_spin_lock(pthread_spinlock_t*)`
`pthread_spin_trylock(pthread_spinlock_t*)`
`pthread_spin_unlock(pthread_spinlock_t*)`
- **Not** related to condition variables because threads do not wait and are not woken up!

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Semaphores (see PSS)

- Special counter
 - inc & dec atomic
 - no access to its value
 - wait for counter > 0 & dec
 - inc & signal a blocked thread/process.
- Initial counter
 - if == 0, useful for synchronizing.
 - if == *n* (> 0), useful for allowing at most *n* threads/processes in a critical section.

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Semaphores

```
#include <semaphore.h>
```

- Calls
 - `sem_init(sem_t*, int, unsigned int value)`
`sem_destroy(sem_t*)`
 - `sem_wait(sem_t*)`
`sem_trywait(sem_t*)`
`sem_timedwait(sem_t*, const struct timespec*)`
 - `sem_post(sem_t*)`

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Producer-Consumers

```

Initial value: 0
Semaphore S
  |
  |
  v
Producer do
...
sem_wait(S)
write(&data)
sem_post(M)
loop

Shared data

Consumer do
...
sem_wait(M);
if hasData(&data)
read(&data)
sem_post(M)
else
sem_post(S)
loop


Initial value: 1
Semaphore M
  
```

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Application to our Model-Checker

- Spin-locks on queues.
- Semaphores for synchronization:
 - producer-consumer scheme!
 - Careful with races to detect termination.
 - Atomic: modify the queue & keep track of blocked threads with a flag.

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 **Advanced: Futex**

- Futex: Fast userspace locking system call.
 - Wait for a value at a given address to change.
 - Wake up anyone waiting on an address.
- Low-level call usually used to implement locks.
- Minix specific, available under Linux.

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