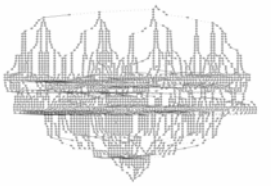
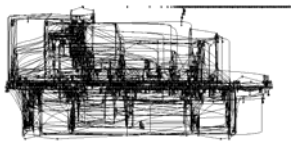
Applications of BDDs

1. **BDDs – short review**
2. **BDDs and Sudoku**
3. **BDDs and verification**
4. **SAT-solving versus BDDs**





Binary Decision Diagrams

[Randal Bryant'86]

A short review

ROBDDs formally

A *Binary Decision Diagram* is a rooted, directed, acyclic graph (V, E). V contains (up to) two *terminal* vertices, $0, 1 \in V$. $v \in V \setminus \{0, 1\}$ are *non-terminal* and has attributes $var(v)$, and $low(v), high(v) \in V$.

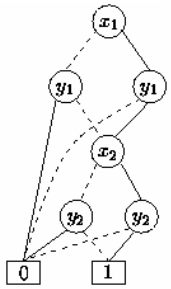
A BDD is *ordered* if on all paths from the root the variables respect a given total order.

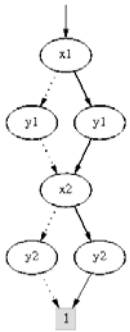
A BDD is *reduced* if for all non-terminal vertices u, v,

1) $low(u) \neq high(u)$
2) $low(u) = low(v), high(u) = high(v), var(u) = var(v)$ implies $u = v$

Reduced Ordered Binary Decision Diagrams

$(x_1 \Leftrightarrow y_1) \wedge (x_2 \Leftrightarrow y_2)$



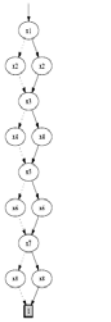


Iben

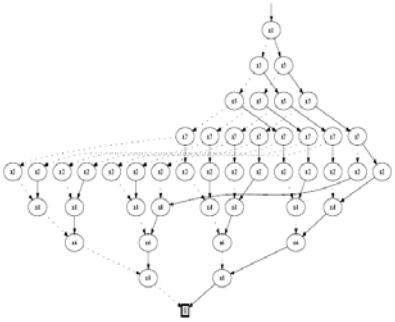
Edges to 0 implicit

Ordering DOES matter

$(x_1 \Leftrightarrow x_2) \wedge (x_3 \Leftrightarrow x_4) \wedge (x_5 \Leftrightarrow x_6) \wedge (x_7 \Leftrightarrow x_8)$



$x_1 < x_2 < \cdots < x_8$



$x_1 < x_3 < x_5 < x_7 < x_2 < x_4 < x_6 < x_8$

Canonicity of ROBDDs

$t_0 = 0$
$t_1 = 1$
$t_u = x \rightarrow t_h, t_l$, if u is a node (x, l, h)

Lemma 1 (Canonicity lemma) *For any function $f : \mathbb{B}^n \rightarrow \mathbb{B}$ there is exactly one ROBDD b with variables $x_1 < x_2 < \ldots < x_n$ such that*

$$t_b[v_1/x_1, \ldots, v_n/x_n] = f(v_1, \ldots, v_n)$$

for all $(v_1, \ldots, v_n) \in \mathbb{B}^n$.

Consequences: b is a tautology, if and only if, $b = \boxed{1}$
b is satisfiable, if and only if, $b \neq \boxed{0}$

Build

Let t be a boolean expression and $x_1 < x_2 < \cdots < x_n$.
Build$(t, 1)$ builds a corresponding ROBDD and returns its root.

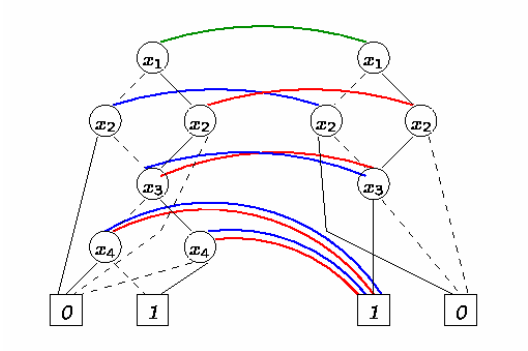
Build(t, i): Node =
if $i > n$ **then**
  if t is true **then return** 0 **else return** 1
else
  low := Build$(t[0/x_i], i+1)$
  $high$:= Build$(t[1/x_i], i+1)$
  var := i
  return Makenode$(var, low, high)$
end if

Complexity ??

APPLY operation

```
Apply(op, b1, b2)
4:  function app(u1, u2) =
6:          if u1 ∈ {0,1} and u2 ∈ {0,1} then res ← op(u1, u2)
7:          else if u1 ∈ {0,1} and u2 ≥ 2 then
8:                  res ← makenode(var(u2), app(u1, low(u2)), app(u1, high(u2))
9:          else if u1 ≥ 2 and u2 ∈ {0,1} then
10:                 res ← makenode(var(u1), app(low(u1), u2), app(high(u1), u2))
11:         else if var(u1) = var(u2) then
12:                 res ← makenode(var(u1), app(low(u1), low(u2)),
                                           app(high(u1), high(u2)))
13:         else if var(u1) < var(u2) then
14:                 res ← makenode(var(u1), app(low(u1), u2), app(high(u1), u2))
15:         else (* var(u1) > var(u2) *)
16:                 res ← makenode(var(u2), app(u1, low(u2)), app(u1, high(u2)))
18:         return res
20:
21: b.root ← app(b1.root, b2.root)
22: return b
```


APPLY example



APPLY operation
with dynamic programming

```
Apply(op, b1, b2)
4:  function app(u1, u2) =
5:          if G(u1, u2) ≠ empty then return G(u1, u2)
6:          else if u1 ∈ {0,1} and u2 ∈ {0,1} then res ← op(u1, u2)
7:                  else if u1 ∈ {0,1} and u2 ≥ 2 then
8:                          res ← makenode(var(u2), app(u1, low(u2)), app(u1, high(u2))
9:                  else if u1 ≥ 2 and u2 ∈ {0,1} then
10:                         res ← makenode(var(u1), app(low(u1), u2), app(high(u1), u2))
11:                 else if var(u1) = var(u2) then
12:                         res ← makenode(var(u1), app(low(u1), low(u2)),
                                                   app(high(u1), high(u2)))
13:                 else if var(u1) < var(u2) then
14:                         res ← makenode(var(u1), app(low(u1), u2), app(high(u1), u2))
15:                 else (* var(u1) > var(u2) *)
16:                         res ← makenode(var(u2), app(u1, low(u2)), app(u1, high(u2)))
17:                 G(u1, u2) ← res
18:                 return res
20:
2:  forall i ≤ max(b1), j ≤ max(b2) : G(i, j) ← empty
21: b.root ← app(b1.root, b2.root)
22: return b
```

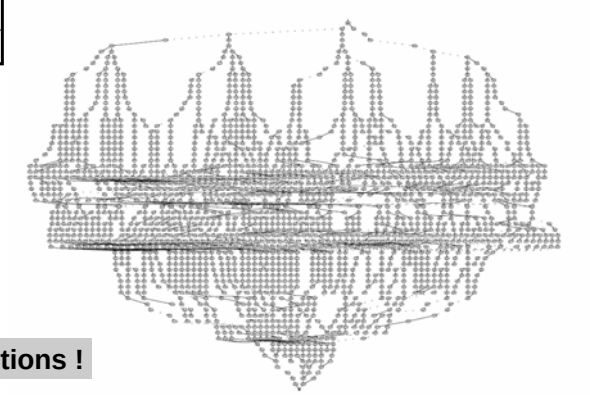

Other operations

Restrict	$b[v/x]$
Size (satcount)	$size(b) = \lvert\{\rho \mid b[\rho] = 1\}\rvert$
Anysat	$anysat(b) = \rho$, for some ρ with $b[\rho] = 1$
Allsat	$allsat(b) = \{\rho \mid b[\rho] = 1\}$
Compose	$compose(b, x, b') = b[x/b']$
Existential quantification	$\exists x.b = b[x/0] \vee b[x/1]$

Constraint Solving & Analysis & IBEN

1			
	2		
		3	
			4

4 x 4 Sudoku



288 solutions !

Encoding

	1	2	3	4
1	1			
2		2		
3			3	
4				4

Boolean variables $x_{i,j,k}$ for all i, j, k ∈ {1,2,3,4}.

Idea:

$x_{i,j,k}$ = 1 if the number k is in position (i,j) in the solution
0 otherwise

$x_{2,2,2}$ =1
$x_{4,4,4}$ =1
$x_{2,2,1}$ =0

Constraints

	1	2	3	4
1	1			
2		2		
3			3	
4				4

Precisely one value in each position i, j:
$x_{1,j,1} + x_{i,j,2} + x_{i,j,3} + x_{i,j,4} = 1$ for each i, j

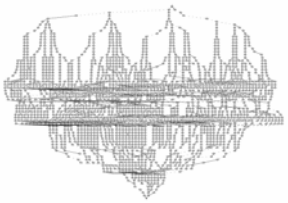
Each value k appears in each row i exactly ones:
$x_{i,1,k} + x_{i,2,k} + x_{i,3,k} + x_{i,4,k} = 1$ for each i, k

Each value k appears in each colomn j exactly ones:
$x_{1,j,k} + x_{2,j,k} + x_{3,j,k} + x_{4,j,k} = 1$ for each j, k

Each value k appears in each 2x2 box exactly ones:
$x_{1,1,k} + x_{1,2,k} + x_{2,1,k} + x_{2,2,k} = 1$ (e.g.)

Solving Sudoku

	1	2	3	4
1				
2				
3				
4				



	1	2	3	4
1	1			
2		2		
3			3	
4				4

In IBEN

4 x4 Sudoku
IBEN-demo

An important puzzle

Per, Kristian, Ole and Jens are to hold an Xmas-party.

Unfortunately, they are almost out of money which severely limits the amount of beer at the party.

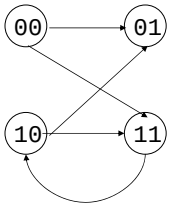
In fact, they have to make do with 1 Tuborg, 1 Carlsberg, 1 Xmas (a Danish Xmas beer), and one Carls Special.

However, the four guys have individual requirements which must be fulfilled at all costs. In particular,
Per only drinks Tuborg and Carlsberg;
Kristian only drinks Carlsberg and Xmas;
Ole essentially drinks everything except Xmas,
and Jens can only drink Carlsberg
and Carls Special.
Is it possible to plan the party drinking so that they all get something to drink?

ROBDDs and Verification

[…,McMillan'90,…..,VVS'97]

ROBDD encoding of transition system

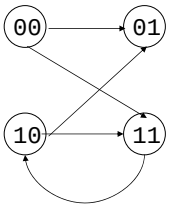


Encoding of states using binary variables (here `x1` and `x2`).

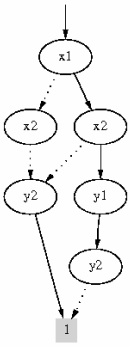
Encoding of transition relation using source and target variables (here `x1`, `x2`, `y1`, and `y2`)

```
Trans(x1,x2,y1,y2):=
     !x1 & !x2 & !y1 & y2
   + !x1 & !x2 &  y1 & y2
   +  x1 & !x2 & !y1 & y2
   +  x1 & !x2 &  y1 & y2
   +  x1 &  x2 &  y1 & !y2;
```


ROBDD representation (cont.)

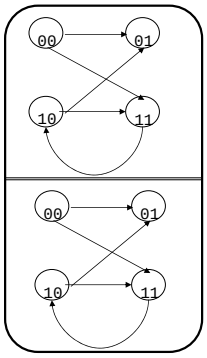



```
Trans(x1,x2,y1,y2):=
     !x1 & !x2 & !y1 & y2
   + !x1 & !x2 &  y1 & y2
   +  x1 & !x2 & !y1 & y2
   +  x1 & !x2 &  y1 & y2
   +  x1 &  x2 &  y1 & !y2;
```



ROBDD for parallel composition

`ATrans(x,y)`



`BTrans(u,v)`

Asynchronous composition

```
Trans(x,y,u,v) =
      (ATrans(x,y) & v=u)
    + (BTrans(u,v) & y=x)
```

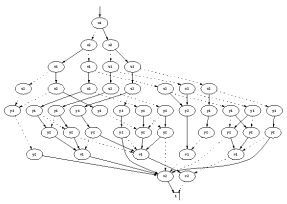
Synchronous composition

```
Trans(x,y,u,v) =
    (ATrans(x,y) & BTrans(u,v))
```

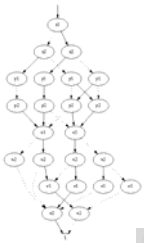
Which ordering to choose?

Ordering?

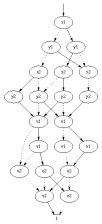
45 nodes
x1,x2,u1,u2, y1,y2 ,v1,v2



23 nodes
x1,x2,y1,y2,u1,u2,v1,v2



20 nodes
x1,y1,x2,y2,u1,v1,u2,v2

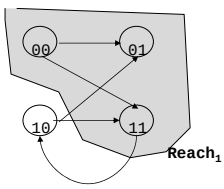


Polynomial size BDDs guaranteed in size of argument BDDs [Enders,Filkorn, Taubner'91]

Reachable States

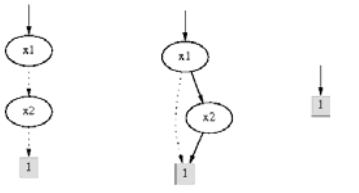
```
Reach(x) := Init(x);
REPEAT
   Old(x) := Reach(x);
   New(y) := Exists x.(Reach(x) & Trans(x,y));
   Reach(x) := Old(x) + New(x)
UNTIL Old(x) = Reach(x)
```

Relational Product: May be constructed without building intermediate (often large) &-BDD.



$Reach_1$

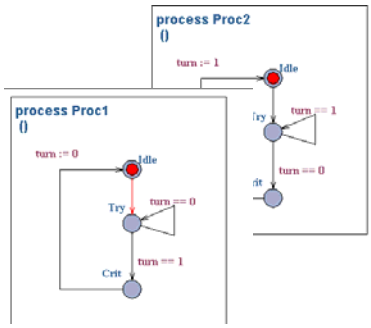
$Reach_0$ $Reach_1$ $Reach_2$



A MUTEX Algorithm

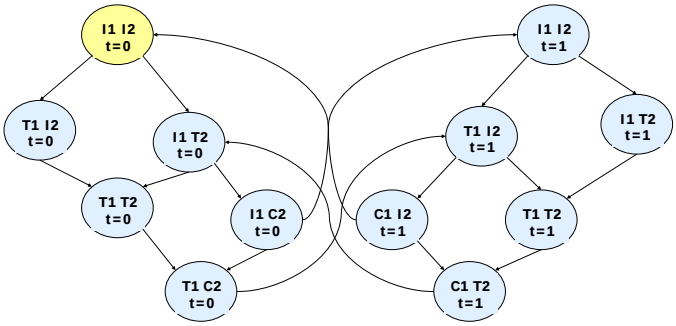
Clarke & Emerson

```
P1 :: while True do
      T1 : wait(turn=1)
      C1 : turn:=0
      endwhile
||
P2 :: while True do
      T2 : wait(turn=0)
      C2 : turn:=1
      endwhile
```



Mutual Exclusion Program

Global Transition System



A MUTEX Algorithm

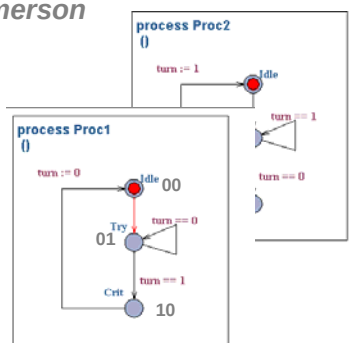
Clarke & Emerson

```
vars x1 x2;
vars y1 y2;
vars u1 u2;
vars v1 v2;
vars t s;

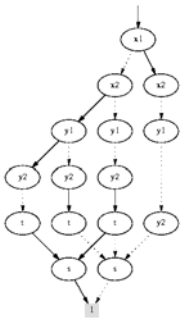
ATrans := (!x1 & !x2 & !y1 & y2 & (s=t))
        + (!x1 & x2 & !y1 & y2 & !t & !s)
        + (!x1 & x2 & y1 & !y2 & t & s)
        + (x1 & !x2 & !y1 & !y2 & !s);

BTrans := (!u1 & !u2 & !v1 & v2 & (s=t))
        + (!u1 & u2 & !v1 & v2 & t & s)
        + (!u1 & u2 & v1 & !v2 & !t & !s)
        + (u1 & !u2 & !v1 & !v2 & s);

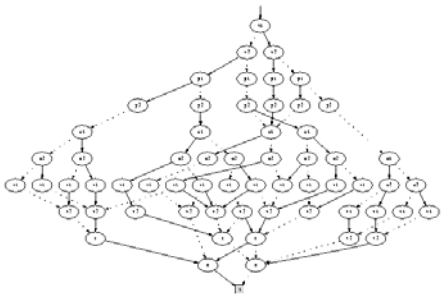
TT := (ATrans & (u1=v1) & (u2=v2))
    + (BTrans & (x1=y1) & (x2=y2));
```



BDDs for Transition Relations



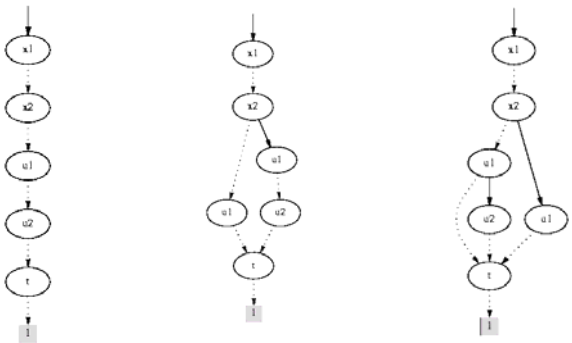
ATrans



TT

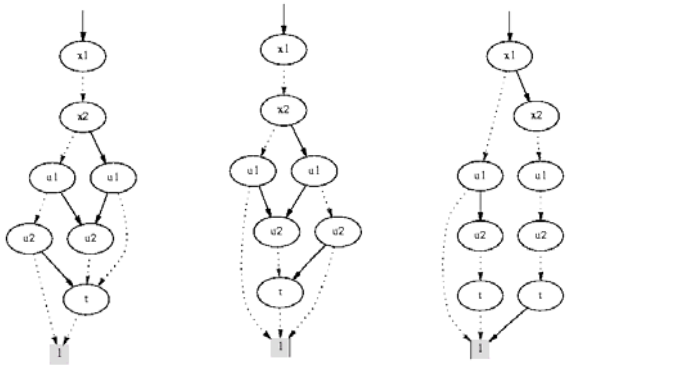
Reachable States

```
Reach(x) := Init(x);
REPEAT
   Old(x) := Reach(x);
   New(y) := Exists x.(Reach(x) & Trans(x,y));
   Reach(x) := Old(x) + New(x)
UNTIL Old(x) = Reach(x)
```



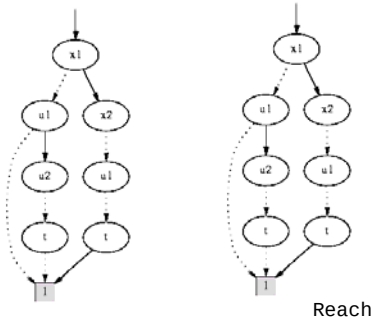
Reachable States

```
Reach(x) := Init(x);
REPEAT
   Old(x) := Reach(x);
   New(y) := Exists x.(Reach(x) & Trans(x,y));
   Reach(x) := Old(x) + New(x)
UNTIL Old(x) = Reach(x)
```



Reachable States

```
Reach(x) := Init(x);
REPEAT
   Old(x) := Reach(x);
   New(y) := Exists x.(Reach(x) & Trans(x,y));
   Reach(x) := Old(x) + New(x)
UNTIL Old(x) = Reach(x)
```

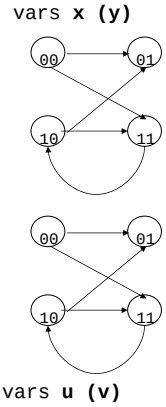


Reach

MUTEX ?

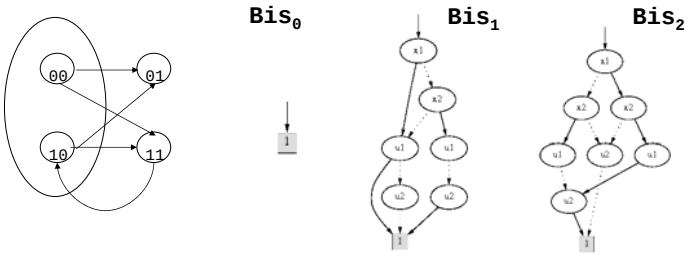
```
Reach &
   x1 & !x2 &
   u1 & !u2
```


Bisimulation



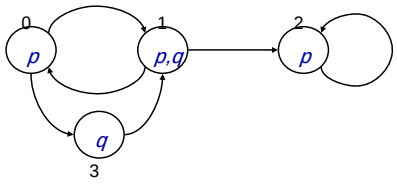

```
Bis(x,u):= 1;
REPEAT
   Old(x,u) := Bis(x,u);
   Bis(x,u) :=
      Forall y. Trans(x,y) =>
         (Exists v. Trans(u,v) & Bis(y,v))
      &
      Forall v. Trans(u,v) =>
         (Exists y. Trans(x,y) & Bis(y,v));
UNTIL Bis(x,u)=Old(x,u)
```


Bisimulation (cont.)



3 equivalence classes
= 6 pairs in final bisimulation

Model Checking



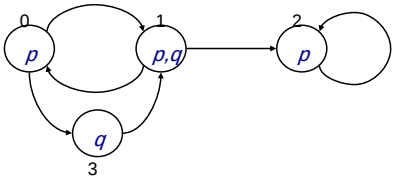

```
vars x1 x2;
vars y1 y2;

Trans(x1,x2,y1,y2) :=
   !x1 & !x2 & !y1 & y2
 + !x1 & !x2 & y1 & y2
 + ……  ;

P(x1,x2) := !x1 & !x2
   + !x1 & x2
   + x1 & !x2;

Q(x1,x2) := …… ;
```

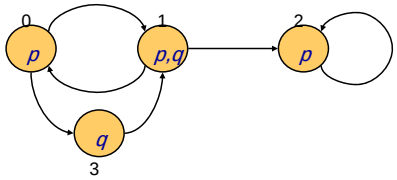

Model Checking



<>P

```
Exists y1,y2.
   Trans(x1,x2,y1,y2) &
   P(y1,y2);
```


Model Checking



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```
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   Trans(x1,x2,y1,y2) &
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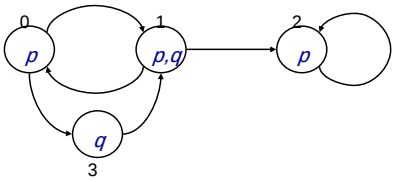
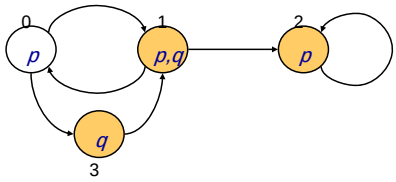
Model Checking



[]P

```
Forall y1,y2.
   Trans(x1,x2,y1,y2) =>
   P(y1,y2);
```


Model Checking



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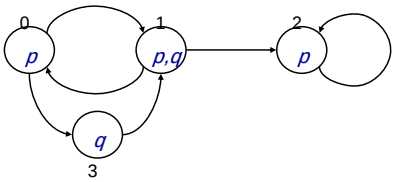
```
Forall y1,y2.
   Trans(x1,x2,y1,y2) =>
   P(y1,y2);
```


Model Checking

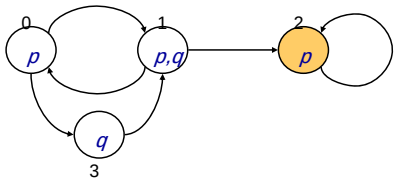


ALWAYS P

max fixpoint

```
A(x1,x2) =
  P(x1,x2) &
  Forall y1,y2.
     Trans(x1,x2,y1,y2) =>
     A(y1,y2);
```


Model Checking

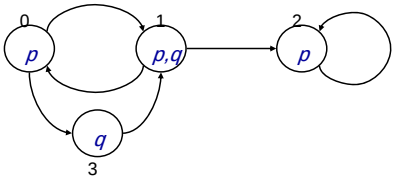


ALWAYS P

max fixpoint

```
A(x1,x2) =
  P(X1,x2) &
  Forall y1,y2.
     Trans(x1,x2,y1,y2) =>
     A(y1,y2);
```


Model Checking

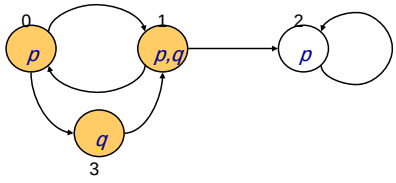


P UNTIL Q

min fixpoint

```
U(x1,x2) =
  Q(x1,x2) +
  { P(x1,x2) &
     Forall y1,y2.
     Trans(x1,x2,y1,y2) =>
       U(y1,y2)
   };
```


Model Checking



P UNTIL Q

min fixpoint

```
U(x1,x2) =
  Q(X1,x2) +
  { P(x1,x2) &
     Forall y1,y2.
     Trans(x1,x2,y1,y2) =>
       U(y1,y2)
   };
```

Partitioned Transition Relation

Relational Product LARGE

```
Exists yv. (T(xu,yv) & S(yv))
```

Asynchronous

```
T(xu,yv) =
      (ATrans(x,y) & v=u)
    + (BTrans(u,v) & y=x)
```

```
Exists y. ATrans(x,y) & S(yu)
+
Exists v. BTrans(u,v) & S(xv)
```

Synchronous

```
T(xu,yv) =
     ATrans(x,y)
   & BTrans(u,v)
```

```
Exists yv.
  Atrans(x,y)
& Btrans(u,v) & S(yv)
```

```
Exists y.
  Atrans(x,y)
& (Exists v. Btrans(u,v) & S(yv))
```

visualSTATE

CIT project VVS (w DTU)

Beologic's Products: salesPLUS visualSTATE

1980-95: *Independent division of B&O*
1995- : *Independent company*
B&O, 2M Invest,
Danish Municipal Pension Ins. Fund
1998: *BAAN*
2000: *IAR Systems A/S*

- *Embedded Systems*
- *Simple Model*
- *Verification of Std. Checks*
- *Explicit Representation (**STATEEXPLOSION**)*
- *Code Generation*

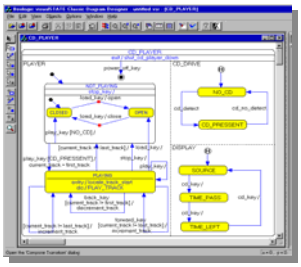
Customers

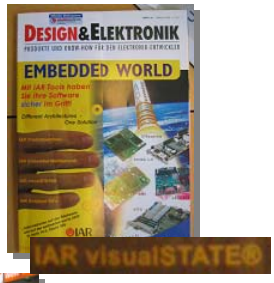
ABB
B&O
Daimler-Benz
Ericson DIAX
ESA/ESTEC
FORD
Grundfos
LEGO
PBS
Siemens ……. (approx. 200)

Verification Problems:
- 1.400 components
- 10^{400} states

Our techniques has reduced verification by an order of magnitude (from 14 days to 6 sec)

visualSTATE



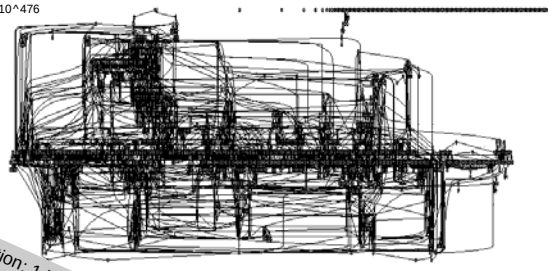


Control Programs

A Train Simulator, visualSTATE (VVS)

1421 machines
11102 transitions
2981 inputs
2667 outputs
3204 local states
Declare state sp.: 10^476

BUGS ?



"Ideal" presentation: 1 bit/state will clearly NOT work!

Experimental Breakthroughs

Patented

System	Mach.	State Space		Checks	Visual	St-of-Art		ComBack	
		Declared	Reach		ST	Sec	MB	Sec	MB
VCR	7	10^5	1279	50	<1	<1	6	<1	7
JVC	8	10^4	352	22	<1	<1	6	<1	6
HI-FI	9	10^7	1416384	120	1200	1.0	6	3.9	6
Motor	12	10^7	34560	123	32	<1	6	2,0	
AVS	12	10^7	1438416	173	3780	6.7	6	5.7	6
Video	13	10^8	1219440	122	---	1.1	6	1.5	6
Car	20	10^11	9.2 10^9	83	---	3.8	9	1.8	6
N6	14	10^10	6399552	443	---	32.3	7	218	6
N5	25	10^12	5.0 10^10	269	---	56.2	7	9.1	6
N4	23	10^13	3.7 10^8	132	---	622	7	6.3	6
Train1	**373**	**10^136**	**---**	**1335**	**---**	**---**	**---**	**25.9**	**6**
Train2	**1421**	**10^476**	**---**	**4708**	**---**	**---**	**---**	**739**	**11**

Machine: 166 MHz Pentium PC with 32 MB RAM

---: Out of memory, or did not terminate after 3 hours.

Experimental Breakthroughs

Patented

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		Declared	Reach		ST			Sec	MB
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JVC	8	10^4	352					<1	6
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Motor	12	10^7					o	2,0	
AVS	12	10^7					6	5.7	6
Video	13	10				1.1	6	1.5	6
Car	20				---	3.8	9	1.8	6
N6				443	---	32.3	7	218	6
N5	25			269	---	56.2	7	9.1	6
N4	23		10^8	132	---	622	7	6.3	6
Train1			**---**	**1335**	**---**	**---**	**---**	**25.9**	**6**
Train2	14		**---**	**4708**	**---**	**---**	**---**	**739**	**11**

Our technique have reduced verification time by several orders of magnitude (eg. From 14 days to 6 sec)

Machine: 166 MHz Penti PC with 32 MB RAM

---: Out of memory, or did not terminate after 3 hours.

SAT-solving

Davis-Putnam

ϕ in CNF, i.e. $\phi = (l_{1,1} \vee \ldots \vee l_{1,n1}) \wedge \ldots \wedge (l_{k,1} \vee \ldots \vee l_{k,nk})$

where $l_{i,j}$ is literal.

clause

```
Function SAT(ϕ)
/unit propagation/
Repeat
  For each unit clause (l) in ϕ
  do  delete from ϕ any clause containing l
      delete !l from any clause in ϕ
  if   ϕ is empty   return  TRUE
  if   ϕ  contains the empty clause   return  FALSE
Until no further change
/splitting/
Choose literal  l  occuring in ϕ   (from smallest clause)
if   SAT(ϕ ∧ (l))  then return  TRUE
else if SAT(ϕ∧ (!l)) then return TRUE
      else  return  FALSE
```

THEOREM

SAT(ϕ)=TRUE
iff
ϕ is satisfiable
-
SAT(ϕ)=FALSE
iff
ϕ is unsatisfiable
-
SAT always terminates

SAT-Solving

- A very active research area with phenomenal performance improvements, e.g.:
  - Analysis of Vital Processor Interlockings (The Netherlands): formulas of size 120K
  - Stålmarks method (PROVER → Trusted Logic) – based on DP plus learning
  - HeerHugo
  - CHAFF – utilizing cash
- Work on SAT-solving for Timed Propositional logic:

  $\phi ::= a \mid x\text{-}y \leq m \mid \neg \phi \mid \phi \wedge \phi$
- See Bulleting of EATCS, February 2005.